

DSC 190

DATA STRUCTURES & ALGORITHMS

Lecture 15 | Part 1

Today's Lecture

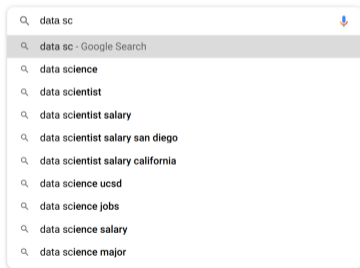
String Data Structures

- ▶ One of the themes of this quarter:
- ▶ If you're doing something once, use an algorithm.
- ▶ If you're doing it over and over, use an appropriate data structure.

String Data Structures

- ▶ Over the next two lectures, we'll look at data structures for strings.
- ▶ Today: **tries** for efficient repeated **prefix queries**.

Autocompletion



DSC 190

DATA STRUCTURES & ALGORITHMS

Lecture 15 | Part 2

Tries

Trie

- ▶ A data structure for storing strings.
- ▶ Pronounced “try”, short for “re**trieval**”.
- ▶ Supports fast **prefix query** and **membership query**.

Prefixes

- ▶ A **prefix** p of a string s is a contiguous slice of the form $s[0:t]$, for some t .
- ▶ Examples:
 - ▶ "test" is a prefix of "testing"
 - ▶ "te" is a prefix of "testing"
 - ▶ "sa" is a prefix of "san diego"
 - ▶ "di" is **not** a prefix of "san diego"

Prefix Query

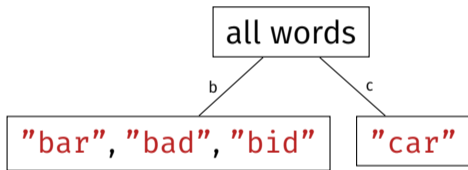
- ▶ **Given:** a collection of n strings and a prefix, p .
- ▶ **Find:** all strings in the collection for which p is a prefix.
- ▶ Example:
 - ▶ "bar", "bad", "bid", "car"
 - ▶ $p = "ba"$

Brute Force

- ▶ Loop over each of n strings, compare against prefix p .
- ▶ Worst-case time: $\Theta(n \cdot |p|)$

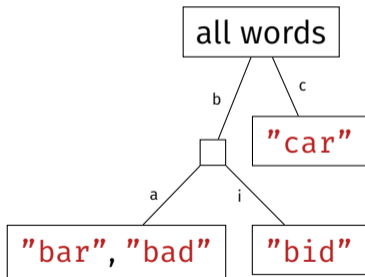
Trie: Motivation

"bar", "bad", "bid", "car"



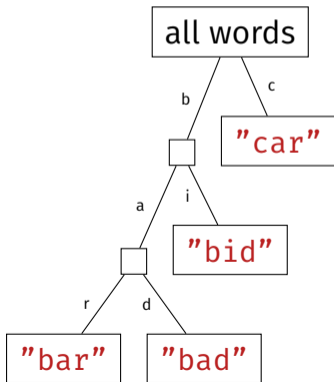
Trie: Motivation

"bar", "bad", "bid", "car"

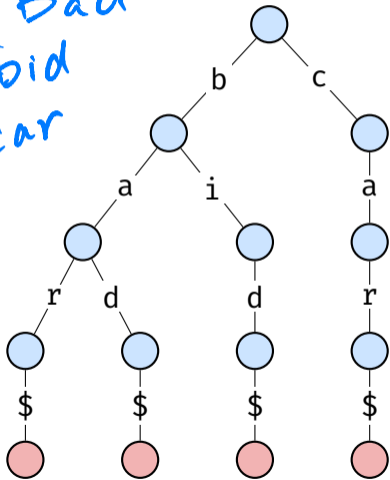


Trie: Motivation

"bar", "bad", "bid", "car"



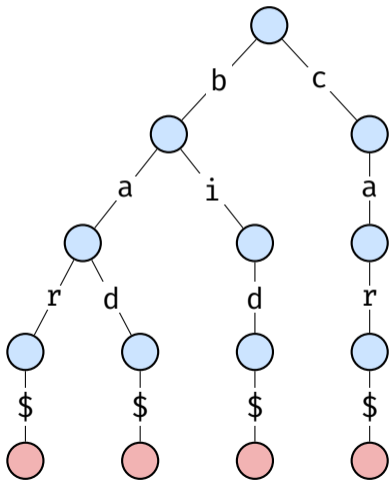
bar
bad
bid
car



Tries

- ▶ Internal nodes represent prefixes.
- ▶ Leaf nodes represent full words.
- ▶ Edges are characters.
- ▶ Words are encoded as paths.

Sentinels



- ▶ \$ is a **sentinel**.
- ▶ It is different from the dollar sign character.
- ▶ It marks the end of a word.
- ▶ Used to show that **"bar"** in trie, but **"ba"** not.

DSC 190

DATA STRUCTURES & ALGORITHMS

Lecture 15 | Part 3

Implementing Tries

S = "\$"

Representation

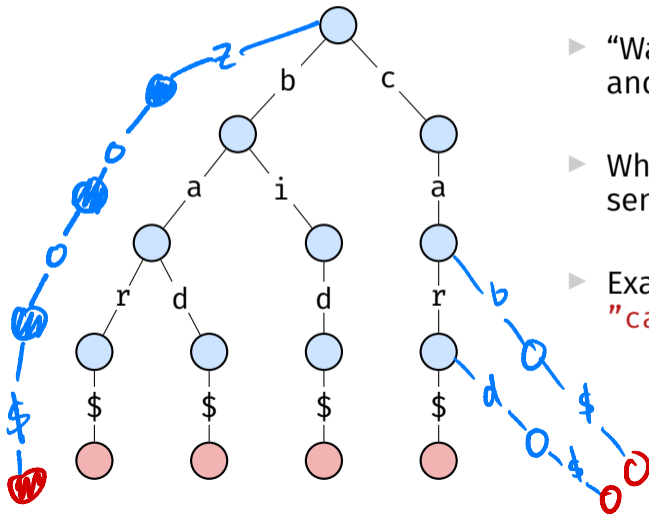
- ▶ Each node has a hash table / array mapping characters to a child nodes.
- ▶ Sentinel represented with a singleton object?

```
END_OF_STRING = object()
```

```
class TrieNode:
```

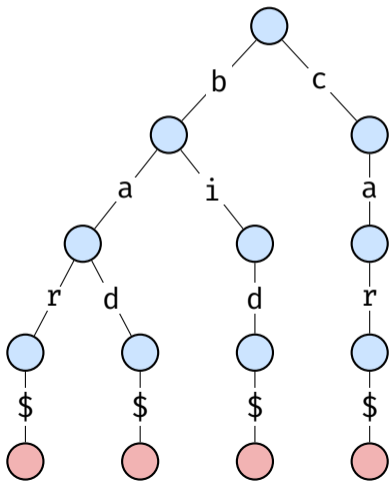
```
    def __init__(self):  
        self.children = {}
```


Insertion



- ▶ “Walk” down tree, creating edges and nodes as necessary.
- ▶ When no more letters left, add sentinel.
- ▶ Example: insert “cab”, “card”, “zoo”

Insertion (Recursive)



- ▶ Suppose we `.insert(s)` on root node.
- ▶ If `s[0]` not in `self.children`, create a new node.
- ▶ Otherwise, let child be `self.children[s[0]]`.
- ▶ Recursively insert `s[1:]` into child.

```
def insert(self, s, start=0, stop=None):
    """Insert s[start:stop] into the trie."""
    if stop is None:
        stop = len(s)

    if start >= stop:
        self.children[END_OF_STRING] = TrieNode()
        return

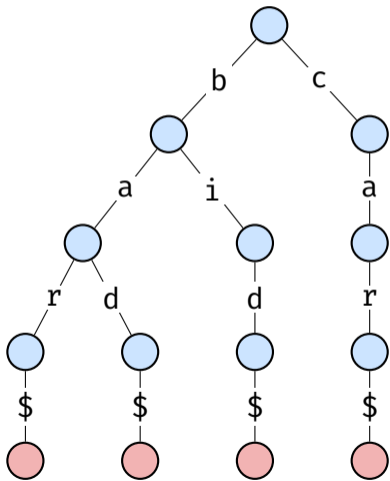
    if s[start] not in self.children:
        self.children[s[start]] = TrieNode()

    child = self.children[s[start]]
    child.insert(s, start + 1, stop)
```

Insertion Time Complexity

- ▶ $\Theta(|w|)$ time, where w is the string inserted.
- ▶ No matter how many elements in trie!

Walk



- ▶ Useful operation.
- ▶ Given a prefix, “walk” down tree.
- ▶ If we “fall off”, raise error.
- ▶ Otherwise, return last node seen.
- ▶ Examples: **“ba”**, **“bo”**

```
def walk(self, s, start=0, stop=None):
    """Walk the trie following s[start:stop].
    Raises ValueError if falls off tree.
    Returns last node encountered otherwise."""
    if stop is None:
        stop = len(s)

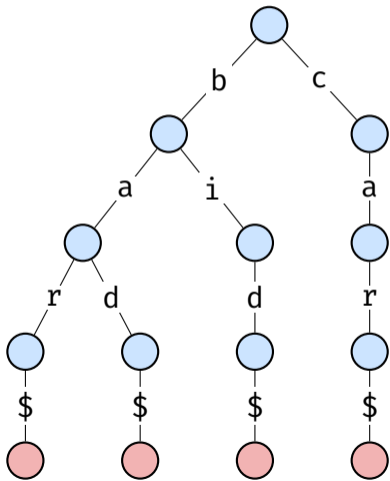
    if start >= stop:
        return self

    if s[start] not in self.children:
        raise ValueError('Fell off tree.')
    else:
        child = self.children[s[start]]
        return child.walk(s, start + 1, stop)
```

Walk Time Complexity

- ▶ Worst-case $\Theta(|p|)$ time, where p is the prefix searched.
- ▶ No matter how many elements in trie!

Membership Query



- ▶ Given p , return True/False if p in collection.
- ▶ “Walk” down tree.
- ▶ If we “fall off”, return False.
- ▶ If not, check that sentinel in children.
- ▶ Examples: **“ba”**, **“bad”**

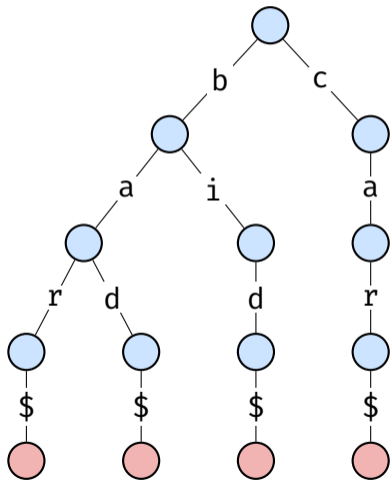

```
def membership_query(self, s, start=0, stop=None):
    """Determine if s[start:stop] is in trie."""
    try:
        node = self.walk(s, start, stop)
    except ValueError:
        return False

    return END_OF_STRING in node.children
```

Membership Query Time Complexity

- ▶ Worst-case $\Theta(|w|)$ time, where w is the prefix searched.
- ▶ No matter how many elements in trie!

Produce



- ▶ Goal: generate all words in subtree.
- ▶ Perform a DFS, keeping track of letters along path.
- ▶ If we find a sentinel, print path.

```
def produce(self, pathchars=''):
    """Generate the words in the trie."""
    for letter, child in self.children.items():
        if letter is END_OF_STRING:
            yield pathchars
        else:
            yield from child.produce(pathchars + letter)
```

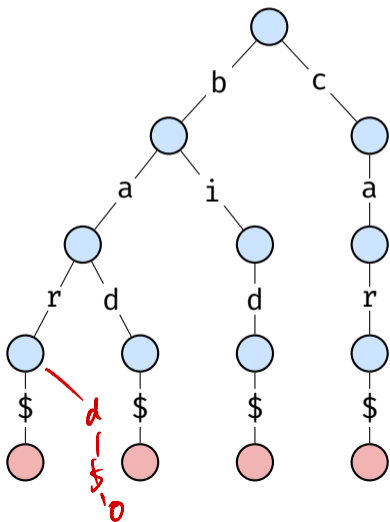
```
def even_numbers():
    x = 0
    while True:
        yield x
        x += 2
```

```
for x in even_numbers():
    print(x)
```

Produce Time Complexity

- ▶ Worst-case $\Theta(\ell)$ time, where ℓ is total length of all strings stored in the trie.
- ▶ If length strings is considered a constant, this is $\Theta(n)$.

Prefix Query (Complete)



- ▶ Given p, return all completions.
- ▶ “Walk” down tree.
- ▶ If we “fall off”, return empty list..
- ▶ If not, produce all nodes in subtree.
- ▶ Examples: **“ba”**, **“bad”**

sum([x**2 for x in range(6)])

```
def complete(self, prefix):  
    try:  
        node = self.walk(prefix)  
    except ValueError:  
        return []  
    return list(node.produce())
```

↓
list(prefix + s for s in node.produce())

Prefix Query Time Complexity

- ▶ Worst-case $\Theta(|p| + \ell_p)$ time, where p is the prefix searched and ℓ_p is the total length of all matches.
- ▶ If length is considered constant, this is $\Theta(|p| + z)$, where z is number of matches.

DSC 190

DATA STRUCTURES & ALGORITHMS

Lecture 15 | Part 4

Demo